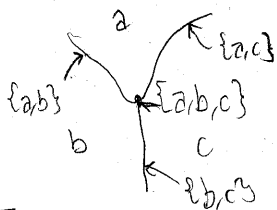


Structure of Rationalizability: applications



Robustness of predictions

• Which type of refinements are robust?

Model: $M = (N, A, u, \Theta \times T, p)$

Solution concept: $\Sigma : M \mapsto \Sigma(M) \subseteq \Delta(A^T)$
strategy profiles

Non-emptiness: Σ is nonempty iff $\Sigma(M) \neq \emptyset$ for each finite M with CPA

Refinements: Σ is a refinement of rationalizability iff $s(t) \in S^\infty[t] \forall t \in T \forall s \in \text{supp}(s)$ and $\forall \sigma \in \Sigma(M)$

Predictions: \exists formula Q with free variable $s: T \rightarrow A$ is a prediction (Σ, M) iff

$Q(s)$ is true $\forall s \in \text{supp}(s), \forall \sigma \in \Sigma(M)$

• This is the stuff that is true under a solution

Metric on T^* : Let d be consistent with the product topology on T^* .

Perturbation: (T', τ) is (ϵ, k) -perturbation of M if

• $T' \subseteq T^*$ is finite, belief-closed, and has a common prior. (T' is a "nice" model)

• $\tau: T \rightarrow T'$ s.t. $\forall t, \forall l \leq k, d(\tau^l(t), t^l) \leq \epsilon$.

Robustness: Q is (ϵ, k) robust if $\forall (\epsilon, k)$ perturbation (T, τ) , $\forall \sigma \in \Sigma(T)$, $\forall s \in \text{supp}(\sigma)$, $Q(s|\sigma, \tau)$ is true.

Q is robust iff it is (ϵ, k) -robust for some $\epsilon > 0$ and $k < \infty$

Thm: \forall nonempty refinement Σ of rationalizability and any finite model M , a prediction Q of (Σ, M) is robust iff Q is a robust prediction of (S^∞, M)

(\Rightarrow) Follows from upper-hemicontinuity of rationalizability.

(\Leftarrow) Follows from definitions

	A	NA
A	θ, θ	$\theta-1, 0$
NA	$0, \theta-1$	$0, 0$

$$\theta = \frac{2}{3}$$

Robustness \Rightarrow A is not robust

CVD, K-M \Rightarrow NA is "robust"

example earlier \Rightarrow NA is not robust

Consider (N, Ω, I, μ)

common prior

want to define a notion of approximate common knowledge under which solution concepts are continuous.

Define $B_i^p(E) = \{\omega : \mu(E|I_i(\omega)) \geq p\}$ p -beliefs

Lemma: $B_i^p(E) \in \mathcal{K}_i$ since this is a property of $I_i(\omega)$ not ω itself.
knowledge field

◦ $E \in \mathcal{K}_i \Rightarrow B_i^p(E) = E, p > 0.$

◦ $E \subseteq F \Rightarrow B_i^p(E) \subseteq B_i^p(F)$ since μ is monotone.
know that you know it
logical omniscience

◦ Consider $E^0 \supseteq E^1 \supseteq E^2 \supseteq \dots$. Then,

$$B_i^p(\bigcap_{n=1}^{\infty} E^n) = \bigcap_{n=1}^{\infty} B_i^p(E^n)$$

◦ $\mu(E | B_i^p(E)) \geq p$

Defn: E is p -evident iff $\forall i, E \subseteq B_i^p(E)$ (everything is $p=0$ -evident)
 (cf: E is evident iff $\forall i, E \subseteq K_i(E)$)

Suppose $\theta \in \{0, 1\}$, and suppose an auctioneer announces price that each person hears with probability $1-p$.

$\theta \backslash s_i$	0	1
0	p	$1-p$
1	$1-p$	p

$\theta_i=0, s_i=0, s_j=0$
 ◦ $E = \{(0, 0, 0)\}$

◦ $P_i[\theta=0 | s_i=0] = \frac{p}{p+(1-p)} = p$

◦ $P_i[\theta=0 | s_i=1] = 1-p$

◦ $B_i^p(E) = \{s_i=0\} \supseteq E$

Let $E = \{\theta=0\}$. Is this p^2 evident?

$B_i^{p^2}(E) = \{s_i=0\} \neq E.$

If $\pi > 1-p$, then E cannot be π evident

$$\text{let } x_i = \theta + \eta_i$$

$$\eta_i \sim U[-\varepsilon, \varepsilon]$$

◦ can we have p -evidence for $p \geq \frac{1}{2}$?

◦ what do p -evident events look like here?

Defn: C is a common p -belief at ω iff $\exists E \ni \omega$ s.t.

i) E is p -evident

ii) $E \subseteq B_i^p(C) \forall i$.

cf: C is common knowledge at ω iff $\exists E \ni \omega$ s.t.

i) E is evident (public)

ii) $E \subseteq C$

$$\circ E = \bigcap_i K_i(E) = \bigcap_i K_i(C)$$

Defn: let $E^p(C) = \bigcap_{n \geq 1} C^n$, where

$$\circ C^0 = C$$

$$\circ C^n = \bigcap_{i=1}^n B_i^p(C^{n-1})$$

Proposition: $\forall C, p,$

1] $E^p(C)$ is p -evident, and $E^p(C) \subseteq B_i^p(C) \forall i$

◦ whenever C occurs, everyone assigns at least probability p to it

2] C is common p -belief at ω iff $\omega \in E^p(C)$

$$\begin{aligned} \text{I] Observe } C^{n+1} &\subseteq C^n \text{ since } C^n \subseteq B_i^P(C^{n-1}) \\ &\Rightarrow B_i^P(C^n) \subseteq B_i^P(B_i^P(C^{n-1})) = B_i^P(C^{n-1}) \\ &\Rightarrow C^{n+1} = \bigcap_i B_i^P(C^n) \subseteq \bigcap_i B_i^P(C^{n-1}) = C^n \end{aligned}$$

$$\text{Next } E^P(C) \subseteq C^{n+1} \subseteq B_i^P(C^n)$$

$$\Rightarrow E^P(C) \subseteq \bigcap_{n \geq 1} B_i^P(C^n)$$

$$\text{Then, } E^P(C) \subseteq \bigcap_{n \geq 1} B_i^P(C^n) = B_i^P\left(\bigcap_{n \geq 1} C^n\right) = B_i^P(E^P(C))$$

$\Rightarrow E^P(C)$ is a p -evident event.

For 2], (\Leftarrow) direction is just I],

$$\circ B_i^P(E) \supseteq E$$

$$\circ B_i^P\left(\bigcap B_i^P(E)\right) \supseteq E \Rightarrow E \subseteq E^P(C)$$

If \underbrace{C} -phabet that posteriors are (π_1, \dots, π_n) , then

$$(\pi_i - \pi_j) \leq 2(1-p)$$

\circ This generalizes the Agreeing to Disagree thm.