

- I Proof by induction carefully
- II Details on problem 3
- III Cooperation in finitely repeated PD.

I Proof by induction

$k \in \mathbb{N} \setminus \{0\}$

$H_k$  - induction hypothesis at stage  $k$

• prove that  $H_1$  holds and that  $H_k \Rightarrow H_{k+1}$ .

### Question 1a

We proceed by induction on the stage of the game  $G$ .

Let  $k$  be the # of non-terminal histories

Denote  $|G| = k$

$H_k$  = any game with  $|G| \leq k$  has a unique strategy that survives iterated elimination of conditionally dominated strategies

To show  $H_1$  holds:

• if  $|G| = 1$ , then it is a 1 person decision problem.

•  $n = 1$  implies  $\exists!$   $u$ -maximizing action  $\Rightarrow H_1$  holds

To show  $\mathcal{H}_k \Rightarrow \mathcal{H}_{k+1}$ :

• Consider  $G$  s.t.  $|G| = k+1$ . Take any SPE of  $G$

• pick  $h$  a final decision node

• at  $h$ ,  $i(h), s_{i(h)}(h)$  maximizes  $i(h)$ 's utility given  $h$

•  $|G| \Rightarrow \exists!$  action  $s_{i(h)}(h)$  consistent with SPE

$\Rightarrow S = (s_{-i} | s_{-i})$  must be SPE of game  $G'$  where  $h$  is turned into final node with payoffs  $u_j(s_{i(h)}(h) | h)$

•  $|G'| = k \quad \square$

Steps: 1] choose something you want to induct on:  $k$

2] Define  $\mathcal{H}_k$

3] Prove  $\mathcal{H}_1$

4] Prove  $\mathcal{H}_k \Rightarrow \mathcal{H}_{k+1}$

E-mail Game by Rubinstein has a good proof by induction in a repeated game.

### Purification results

• mixed eq. in game  $G$  very close to pure eq. in game  $G$ , an elaboration of  $G$

eg: at each decision node, get signal  $u \in [0,1]$   
 $\pi = \pi + 0 \cdot u$  and then action is a function of  $u$ .

Problem 3:  $g(u_2) = \max \{u_1 : (u_1, u_2) \in X\}$

a) From class,  $\exists!$  SPE where

1 proposes  $(x_1, g^{-1}(x_1))$

1 accepts  $(g(x_2), x_2)$

where  $x_1 = g(\delta_2 g^{-1}(\delta_1 x_1))$

here  $g(x_2) = 1 - x_2$

let  $f(x_1) = g(\delta_2 g^{-1}(\delta_1 x_1))$

(\*) More precisely,

$$\begin{aligned} x_1 &= 1 - \delta_2 (1 - \delta_1 x_1) \\ &= 1 - \delta_2 + \delta_1 \delta_2 x_1 \end{aligned}$$

$$\Rightarrow x_1 (1 - \delta_1 \delta_2) = 1 - \delta_2$$

$$\Rightarrow x_1 = \frac{1 - \delta_2}{1 - \delta_1 \delta_2}$$

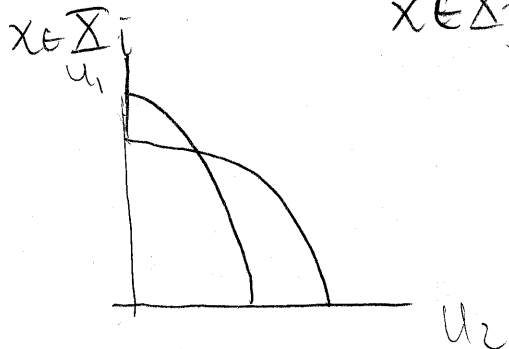
$$x_2 = \frac{1 - \delta_1}{1 - \delta_1 \delta_2}$$

Note that  $x_2 \neq 1 - x_1$

b) Continuity of  $x_1(\delta_1, \delta_2)$  as  $(\delta_1, \delta_2) \rightarrow (1,1)$   
 (Can you extend  $x_1(\cdot, \cdot)$ ?)

(b)  $\underline{X}_1 \neq \underline{X}_2$ 

$$\circ \max_{x \in \underline{X}_i} x_i < \max_{x \in \underline{X}_j} x_i \quad \forall i \neq j$$



$$\circ \text{could have had } \max_{x \in \underline{X}_i} x_i > \max_{x \in \underline{X}_j} x_i \quad \forall i \neq j$$

$$\circ \bar{v}_i = \max_{\{s \text{ survives ICD}\}} u_i(s) \quad \text{does this necessarily exist?}$$

$$\circ \underline{v}_i = \inf_{\{s_j \text{ survives ICD}\}} \sup_{\{s_i \text{ survives ICD}\}} u_i(s)$$

This is a x ante payoffs when 1 starts proposing.

$$\circ \text{Conceptually, } \bar{v}_1 \neq \bar{v}_2, \underline{v}_1 \neq \underline{v}_2$$

Candidate

Right approach (proof strategy)

(b) + (c) together

$$\text{Define: } u_1^{b,0} = 0, \quad u_2^{b,0} = 0$$

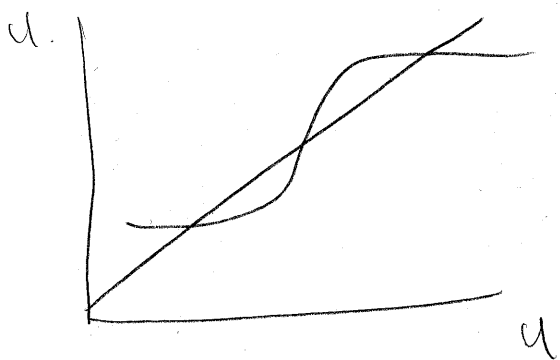
$$u_1^{h,0} = \max_{x \in \underline{X}_1} x_1, \quad u_2^{h,0} = \max_{x \in \underline{X}_2} x_2$$

$$\circ g_1(u_2) = \max \{u_1 \mid (u_1, u_2) \in X_1\}$$

$$\circ g_2(u_1) = \max \{u_2 \mid (u_1, u_2) \in X_2\}$$

Then  $u_1^{k+1} = g_1(g_2 u_2^{k,0})$ , which exists by the weird assumption.

In the end, we get  $u_1^{k,\infty}, u_2^{k,\infty}, u_1^{k+1,\infty}, u_2^{k+1,\infty}$



$f_\varepsilon \rightarrow f$  uniformly over  $K$  as  $\varepsilon \rightarrow 0$

$Q^0 =$  when f.p. of  $f_\varepsilon$  close to f.p. of  $f$

When do extreme fixed points of  $f \circ f_2$  converge to extreme fixed points of  $f \circ f_1$ ?