

Optimal auctions (private iid values)

- Risk neutrality for everybody
- continuum of types
- $\theta_i \sim f_i(\theta_i), F_i(\theta_i)$ (we are temporarily dropping symmetry assumption)
pdf/cdf depend on i
- assume independent private values

Consider a multi-unit auction:

$$(F): \sum_{i=1}^N q_i(\theta) \leq q_0$$

vector of announced types: $\theta = (\theta_1, \dots, \theta_N)$

$$\circ u_i(q, T) = \theta_i v(q) - T$$

What is the set of implementable mechanisms?

- must satisfy (IR) and (IC) in expectation as well as feasibility. (cover others' types)

$$(IC): \pi_i(\theta_i, \theta_{-i}) = \max_{\hat{\theta}_i} \pi_i(\hat{\theta}_i, \theta_{-i})$$

announced type true type

where $\pi_i(\hat{\theta}_i, \theta_{-i}) = \mathbb{E}_{\theta_{-i}} [\theta_i v(q_i(\hat{\theta}_i, \theta_{-i})) - T_i(\hat{\theta}_i, \theta_{-i})]$

true types looking for BNE with truth-telling

Call $T_i(\hat{\theta}_i) = \mathbb{E}_{\theta_{-i}} [T(\hat{\theta}_i, \theta_{-i})]$

It is easy to show that $(IC) \Leftrightarrow (LIC) + (M)$,

where: $(M) \Leftrightarrow \mathbb{E}_{\theta_{-i}} [q_i(\theta_i, \theta_{-i})]$ increasing in θ_i .

This follows from Spence-Mirrlees, since:

$$\frac{\partial}{\partial \theta} \left[\frac{\partial u_i / \partial q}{\partial u_i / \partial (T)} \right] = v'(q) > 0.$$

(LIC): Thus, $\pi_i(\theta_i, \theta_c) = \pi_i(\underline{\theta}, \underline{\theta}) + E_{\theta_i} \left[\int_{\underline{\theta}}^{\theta_i} v(q_i(x, \theta_i)) dx \right]$
 where $\underline{\theta}$ is the lower bound on the support of the distributions.

The participation constraints are:

$$\pi_i(\theta_i, \theta_i) \geq 0 \quad \forall i, \quad \forall \theta_i \quad (IR_{\theta_i})$$

If $(IR_{\underline{\theta}})$ holds, then all (IR_{θ}) hold.

The (LIC) gives us:

$$T_i(\theta_i) = E_{\theta_i} \left[\theta_i v(q_i(\theta)) - \pi_i(\underline{\theta}, \underline{\theta}) - \int_{\underline{\theta}}^{\theta_i} v(q_i(x, \theta_i)) dx \right]$$

$$\text{Then, } T_i^e = \int_{\underline{\theta}}^{\bar{\theta}} E_{\theta_i} \left[\theta_i v(q_i(\theta)) - \pi_i(\underline{\theta}, \underline{\theta}) - \int_{\underline{\theta}}^{\theta_i} v(q_i(x, \theta_i)) dx \right] dF_i(\theta_i) \\ = f_i(\theta_i) d\theta_i$$

Total expected revenue is: $TE = \sum_{i=1}^N T_i^e$

• depends on: rent of lowest type and allocation of the goods.

• This is the revenue equivalence theorem.

Revenue Equivalence Theorem

In the continuum case (with private, independent values and risk neutrality), seller revenue is uniquely determined by $\pi_i(\underline{\theta}, \underline{\theta})$ and $q_i(\underline{\theta})$.

Remark: If standard auctions all have the same $\pi_i(\underline{\theta}, \underline{\theta})$ and $q_i(\underline{\theta})$, they will have the same expected revenue.

◦ Vickrey: people bid their own valuation and winner pays 2nd highest value

◦ ex-post efficient (and $\pi_i(\underline{\theta}, \underline{\theta}) = 0$)

Can use integration by parts to get

$$T_i^E = \int_{\underline{\theta}}^{\bar{\theta}} E_{\theta_i} \left[\theta_i V(q_i(\theta)) - \pi_i(\underline{\theta}, \underline{\theta}) - \frac{V(q_i(\theta_i, \theta_i))}{h_i(\theta_i)} \right] f_i(\theta_i) d\theta_i$$

$$\text{where } h_i(\theta_i) = \frac{f_i(\theta_i)}{1 - F_i(\theta_i)}$$

Optimal Auctions: ($\Rightarrow \pi_i(\underline{\theta}, \underline{\theta}) = 0$)

Assume:

◦ symmetry: $(f_i(\cdot) = f(\cdot))$

◦ $\theta_i v(q) = \begin{cases} \theta_i q & \text{if } q \leq 1 \\ \theta_i & \text{if } q > 1 \end{cases}$

◦ care only about one unit of consumption.

Here, we have:

$$\max_{q_i(\theta)} E_{\theta} \left[\sum_{i=1}^N \left\{ \theta_i \min\{q_i, b\} - \frac{\min\{q_i, b\}}{h(\theta_i)} \right\} \right]$$

$$\text{s.t. } \underbrace{\sum_{i=1}^N q_i(\theta)}_{(F)} \leq q_0 \quad \forall \theta, \text{ and } (M)$$

$$\Leftrightarrow \max_{q_i(\theta)} E_{\theta} \left[\sum_{i=1}^N J(\theta_i) \min\{q_i, b\} \right], \quad J(\theta_i) = \left[\theta_i - \frac{h(\theta_i)}{f(\theta_i)} \right]$$

s.t. - (F) and (M).

Result: Sell to highest $J(\theta_i)$'s provided they have positive $J(\theta_i)$'s. (and $\sum q_i(\theta) \leq q_0$)

• will have (M) if $J(\theta_i) \uparrow$ in θ_i

• sufficient: $\frac{f_i}{1-f_i} \uparrow$ w/ θ_i

The optimal auction is ex-post efficient, possibly with a floor.

interestingly \Leftrightarrow Vickrey, with a floor! (reserve price)

• floor: $J(\hat{\theta}) = 0$ $\hat{\theta}$ is a floor $\Rightarrow \theta < \hat{\theta} \Rightarrow$ do not participate

First-price auction

Look for a Bayesian equilibrium, where bids symmetric

increase strictly with type \Leftrightarrow ex-post efficiency

Maximization problem for each bidder i :

$$\max_{b_i} (\theta_i - b_i) \underbrace{g_i(b_i)}_{\text{probability of winning}}$$

◦ assume two bidders

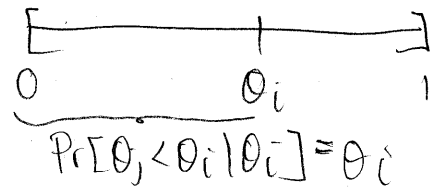
◦ $\theta_i \sim U[0, 1]$ iid.

FOCs:

$$C(b_i): (\theta_i - b_i) g'(b_i) - g(b_i) = 0$$

Symmetry implies $g_i(b_i) = g(b_i) = \theta_i$, since

$$\begin{aligned} \Pr[\text{win}] &= \Pr[\theta_i > \theta_j | \theta_i] = \theta_i \\ &= \Pr[g(\theta_i) > g(\theta_j) | \theta_i] \end{aligned}$$



Thus, the FOCs become:

$$[g(b_i) - b_i] g'(b_i) = g(b_i) \Rightarrow g(b_i) = 2b_i$$

$= \theta_i$

$$\Rightarrow b_i^* = \frac{\theta_i}{2}$$

If values are not independent, the seller can do better.

Cremer-Maclean