

$$x_i \in M[z_i \times B_i]$$

2 interpretations:

• $x_i(z, b)$ - mass of people at (z, b)

• $x_i(z, b) \rightarrow \pi^i(z, b) = \frac{x_i(z, b)}{\sum_{(z', b')} x_i(z', b')}$ lottery

For the quasilinear model, it does not matter which interpretation you use, because associated with $x_i(z, b)$ are money payments $m_i = -pz + P_i(b)$

$\max_{(z', b')} \{v_i(z', b') - pz' + P_i(b')\}$ - you are restricted to
(1) no lotteries

$\max_{(z', b')} \{ \langle v_i, \pi \rangle - \langle (p, P_i), \pi \rangle \}$ (2)
 $\sum v_i(z, b) \pi(z, b)$ weighted average of money payments

(can let $\pi = \delta_{(z, z)}$; (1) is special case of (2))

It turns out there are no gains to lotteries:

Let $\pi = \operatorname{argmax}_{(z', b')} \{v_i(z', b') - pz' + P_i(b')\}$

Suppose we have NTU

• i.e. we insist $\langle (p, P_i), \pi \rangle = 0$

First look at TU economies:

$v_i(q_i)$ concave

$$\max \left\{ \sum_{i=1}^n v_i(q_i) : \sum_{i=1}^n q_i = 0 \right\} = \sum_{i=1}^n v_i(\bar{q}_i)$$

By concavity $\exists p \in \partial v_i(\bar{q}_i) \quad \forall i$

$$\Rightarrow v_i^*(p) = v_i(\bar{q}_i) - p\bar{q}_i$$

$$v_i^*(p) = v_i(\bar{q}_i) - p\bar{q}_i$$

What if $p\bar{q}_i = 0$?

$D_i(p; v_i) = \operatorname{argmax} \{ v_i(q) : pq = 0 \}$
 demand for i given
 utility is v_i

If we can find money transfers s.t. $m_i - pq_i = 0 \quad \forall i$,
 then we are in the NTU economy.

This is sensitive to transformations because of the
 concavity assumption

Suppose $p\bar{q}_i > 0, p\bar{q}_i < 0$. Then v_i

$$(P) \sum_{i=1}^n v_i(\bar{q}_i) = \max_{(i,q)} \left\{ \sum_i v_i(q) x_i(q) : \sum_q x_i(q) = 1 \quad \forall i, \sum_q q x_i(q) = 0 \right\}$$

Imposing the restriction $x_i \in \{0, 1\}$ is
 not binding if v_i is concave

$(i, q) \rightarrow v_i(q)$ payoff
 $x_i(q)$ level of operation

$$(1) \min \left\{ \sum_i 1 \cdot \pi_i : \pi_i + pq \geq v_i(q) \quad \forall (i, q) \right\} = \min_p \left\{ \sum_{i=1}^n v_i^*(p) \right\}$$

$\Leftrightarrow \pi_i \geq v_i(q) - pq$

$\pi_i = v_i^*(p)$

In general, $p q_i \neq 0 \quad \forall i$

The person whose $p q_i \geq 0$ has a better ability to "produce utility from q_i "

$$\text{Now, } \max_{\lambda} \left\{ \sum_{i=1}^n \lambda_i v_i(q_i) : \sum_{i=1}^n q_i = 0 \right\} = \sum_{i=1}^n \lambda_i v_i(q_i(\lambda))$$

We will get some $p(\lambda)$

Looking for a $\lambda \gg 0$ s.t. $p(\lambda) \cdot q_i(\lambda) = 0 \quad \forall i$

$$\text{Let } \Lambda = \left\{ (\lambda_i) : \lambda_i \geq 0, \sum_{i=1}^n \lambda_i = 1 \right\}$$

Does $\exists \lambda \in \Lambda$ with $\lambda \gg 0$ s.t. $p(\lambda) \cdot q_i(\lambda) = 0 \quad \forall i$?

Yes, by a fixed point theorem

$h(c) = \max \{cx : Ax \leq b, x \geq 0\} \Rightarrow$ we can find the set of all solutions for all λ .

Define $\text{dom } v_i = \{q : v_i(q) > -\infty\} = S$

$$I_S(q) = \begin{cases} 0 & \text{if } q \in S \\ -\infty & \text{if } q \notin S \end{cases}$$

$$\text{Let } \Phi_i(\lambda) = \lambda_i + \max \{p(\lambda) q_i(\lambda), 0\}, \quad \Phi = [\Phi_i(\lambda)]$$

$$1 + \sum_j \max \{p(\lambda) q_j(\lambda), 0\}$$

$\Phi : \Lambda \Rightarrow \Lambda$, We will want to apply Kakutani's fixed point theorem.