

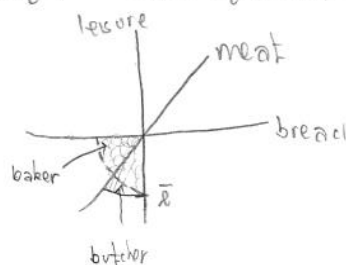
$$E = \{ (U_i, \omega_i, V_i) \}$$

$v_i \in V_i \rightarrow Y_i(v_i)$  "each  $v_i$  defines technology"

$\mathbb{R}^L$  - "super" space

$$\sup_{y_i \in Y_i(v_i)} U_i(\omega_i + y_i + z | v_i) \equiv v_i(z)$$

"best way of attaining the trade 'z'"



leisure	$\begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$	$\begin{pmatrix} 1/2 \\ -1 \\ 0 \end{pmatrix}$	$\begin{pmatrix} 1/2 \\ 0 \\ -1 \end{pmatrix}$
bread			
meat			
	no specialization	make bread (baker)	make meat (butcher)

$$\text{Define } U_i^*(p | v_i) = \sup_z \left\{ \sup_{y_i \in Y_i} U_i(\omega_i + y_i + z | v_i) - p z \right\}$$

$$= \sup_z v_i(z) - p z$$

$$U_i^*(p) = \sup_{v_i} \{ U_i^*(p | v_i) \}$$

PTE for  $E$ :  $p \in \mathbb{R}^L$  ( $z_i$ ) ( $m_i$ ) such that "in PTE price proceed maximization"

i)  $\sum z_i = 0$

ii)  $U_i^*(p) = U_i(\omega_i + y_i + z_i | v_i) - p z_i = v_i(z_i) - p z_i$

Proposition: PTE for  $E$  is efficient.

Pf:  $\forall p, U_i^*(p) \underset{\uparrow v_i}{\geq} U_i^*(p | v_i) \underset{\uparrow z_i}{\geq} v_i(z_i) - p z_i$

By market clearing,  $\sum z_i = 0$

Thus,  $\forall p, \sum U_i^*(p) \geq \sum v_i(z_i)$

Therefore  $\sum U_i^*(p) \geq g(v)$

where  $g(v) = \sup \{ \sum v_i(z_i) : \sum z_i = 0 \}$

By PTE, we have achieved:  $\sum U_i^*(p) = \sum v_i(z_i) - p \sum z_i = \sum v_i(z_i)$

Thus,  $\sum v_i(z_i) \geq g(v) \Rightarrow \sum v_i(z_i) = g(v) \rightarrow PO \quad \square$

Revise the link between prices and maximization:

Instead of  $p \rightarrow v \rightarrow (z_i)$ , we want  $v \rightarrow \tilde{p} \rightarrow (\tilde{z}_i)$

Two-stage game:

Stage 1: Choose  $v$

Stage 2: Payoffs are determined as  $\tilde{\pi}_i(v) = \tilde{\pi}_i(p(v)) = v_i^*(p(v))$

[There are no actions here, so it is effectively one-stage.]

$(\{\tilde{\pi}_i\}, v)$  is a game in normal form  $[v = v_1 \times \dots \times v_n]$

$v \xrightarrow{\text{determines}} \ell(v)$

$$H(v) = \{c : z_{ic} < 0 \text{ for some } v_i (z_i) > -\infty\}$$

$$\mathbb{R}^{\ell(v)} \subset \mathbb{R}^{\ell}$$

Equilibrium:  $\tilde{\pi}_i(v) \geq \tilde{\pi}_i(v_{-i}, v_i')$   $\forall i, \forall v_i'$

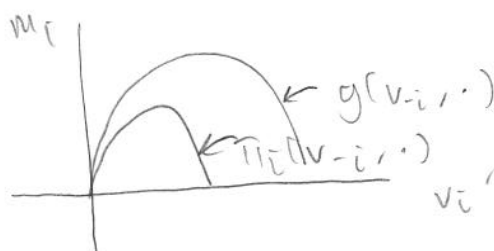
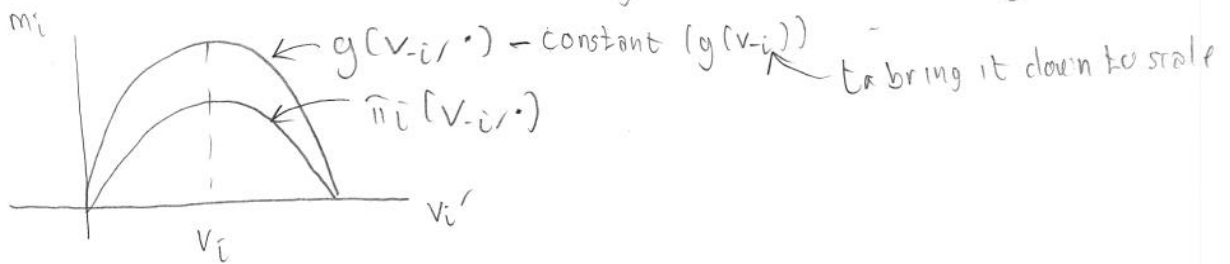
Efficiency wrt  $v \in \mathbb{R}^{\ell(v)}$  (when  $\mathbb{R}^{\ell(v)} = \mathbb{R}^{\ell}$ )

$$\frac{\Delta \tilde{\pi}_i(v)}{\Delta v_i} = \tilde{\pi}_i(v_{-i}, v_i') - \tilde{\pi}_i(v)$$

$$\frac{\Delta g(v)}{\Delta v_i} = g(v, v_i') - g(v)$$

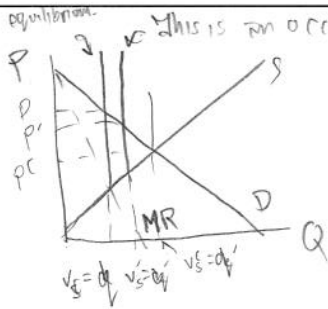
Is it the case that  $\frac{\Delta \tilde{\pi}_i(v)}{\Delta v_i} \frac{\Delta g(v)}{\Delta v_i} \geq 0$

Problem: We have non-efficiency if  $\Delta \tilde{\pi}_i(v) \Delta g(v) < 0$



This is a problem.  
 $\Delta \tilde{\pi}_i(v) \Delta g(v) < 0$

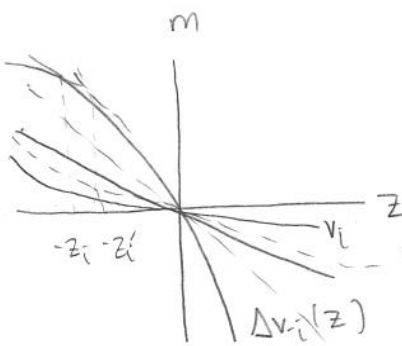
This is what Adam Smith "saw"



This is an occupational choice.

"Can I change my characteristics in a favorable way?"

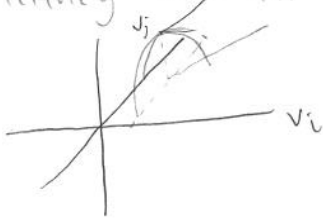
- Misrepresentation of preferences in order to get a more favorable prices



$v_i'$  ← not true preferences (reported)  
- new, more favorable prices for you.

This would not occur under perfect competition. ( $\Delta \pi_i(v) = \Delta g(v)$ )

Efficiency over all  $v$



$$\Delta \pi_i(v) = \Delta g(v) \quad (FA)$$

It is possible that

$$\frac{\Delta g(v)}{\Delta v_i} \leq 0 \quad \forall i, \text{ but } \frac{\Delta g(v)}{\Delta v} > 0$$

$$\tilde{p} \in \mathbb{R}^2$$

← let  $\tilde{p}$  be the price that "precedes maximization"

It would be nice if  $\tilde{p} = (p_c)$

$$\begin{cases} -p_c & c \in H(v_i, v_i') \setminus H(v) \text{ "new commodities that } i \text{ could produce"} \\ p_c & c \in H(v) \\ p_c & c \in H(v_j, v_j') \setminus H(v) \text{ "new commodities that } j \text{ could produce"} \end{cases}$$

This  $\tilde{p}$  is constructed without regards to coordination problem