

Homework due Friday at 12:00pm in the grad. lab.

Consumer:

prices for non-money commodities. $p_m = 1$
 $- d(v, p, w)$ $v: \mathbb{R}^L \rightarrow \mathbb{R} \cup \{-\infty\}$

Utility

$v(z) = -\infty \leftarrow$ not a feasible trade (negative consumption)
 \hookrightarrow we are only interested in $z \ni v(z) > -\infty$

$$d(v, p, w) = \{ (z, m) \mid \underbrace{p \cdot z + m = w}_{\text{market value of trades}}, v(z) + m \geq v(z') + m' \ \forall (z', m') \ni p \cdot z' + m' = w \}$$

Special case: $d(v, p, 0)$ - wealth is zero

$$d(v, p, w) = d(v, p, 0) + (0, w)$$

$$(z^*, m^*) = (z', m') + (0, w)$$

$\Rightarrow z^* = z' \Rightarrow$ No income effects for non-money commodities

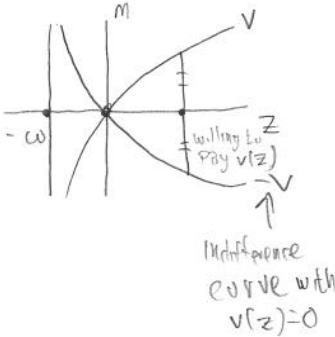


\Rightarrow Wlog, we can assume $w = 0$.

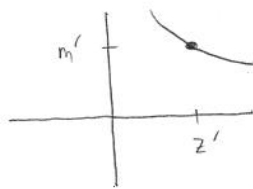
Inverse demand:

$$m = -pz \text{ since } w = 0 \text{ and } pz + m = 0$$

$$d^{-1}(v, z, 0) = \{ p \mid (z, -pz) \in d(v, p, 0) \}$$



$$(z, m) \succeq (z', m') \Leftrightarrow v(z') + m' \leq v(z) + m$$

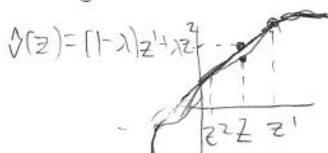


$d^{-1}(v, z, 0)$ can be empty if v is not concave.

Let $\partial v(z) = \{ p \mid v(z) - pz \geq v(z') - pz' \ \forall z' \} \equiv d^{-1}(v, z, 0)$
 \hookrightarrow subdifferential of v at z ,

Define conc $\hat{v} = \hat{v} \ni \hat{v}(z) = \sup \{ \sum \lambda_k v(z_k) \mid \sum \lambda_k z_k = z, \lambda_k \geq 0, \sum \lambda_k = 1 \}$

"concavified utility function" - smallest concave function lying on or above v .
 Clearly, $\hat{v}(z) \geq v(z) \ \forall z$



Proposition: If $p \in \partial v(z)$, $v(z) = \hat{v}(z)$.

↳ If p is in the subdifferential of v at z , then z is not in the non-concave portion of the utility fcn.

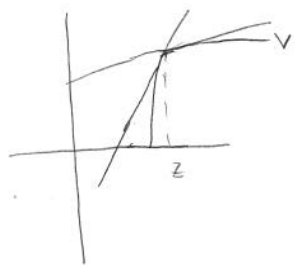
Proof: By contradiction. $v(z) < \hat{v}(z) \Rightarrow p \notin \partial v(z)$. Fill in the details.

Uniqueness of $\partial v(z)$

Suppose v is concave and $\partial v(z) = \{p\}$

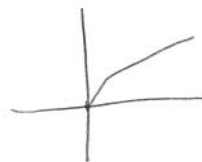
$$\nabla v(z) = \left(\frac{\partial v(z)}{\partial z_1}, \dots, \frac{\partial v(z)}{\partial z_n} \right) = p \Rightarrow \partial v(z) = \nabla v(z)$$

$$\text{[if } l=1, \nabla v(z) = v'(z) = \frac{dv}{dz}(z)\text{]}$$



v is not differentiable at z .

This notion handles boundary problems.



$$\begin{aligned} \text{Indirect utility } (v, p, w) &\equiv \bar{U}(v, p, w) \\ &= v(z) + m \quad (z, m) \in d(v, p, w) \end{aligned}$$

$$\begin{aligned} \text{Indirect utility } (v, p, 0) &\equiv \bar{U}(v, p, 0) \\ &= v(z) + m - w \end{aligned}$$

$$\Rightarrow \bar{U}(v, p, w) - \bar{U}(v, p, 0) = w$$

\Rightarrow Wlog, we can assume $w=0$. this contains the restriction

$$\text{Define } \bar{U}(v, p, 0) = v^*(p) = \sup_z \{ v(z) - pz \}$$

$$\begin{array}{l} v^*(p) \\ \uparrow \\ \text{conjugate of } v \end{array} \quad \begin{array}{l} v(z) \\ v^*(p) \end{array} \quad \begin{array}{l} v: \mathbb{R}^l \rightarrow \mathbb{R} \\ v^*: \mathbb{R}_+^l \rightarrow \mathbb{R} \end{array}$$

Proposition: $v(z) - pz = v^*(p) \Leftrightarrow p \in \partial v(z)$

Pf: This is just a restatement of the definitions.

$$\frac{\partial u_i(x)}{\partial x_k} = \lambda p_k$$

$$\frac{\partial v_i(z)}{\partial z_k} = p_k \quad \text{FOCs for QL utility: } p = \nabla v(z)$$

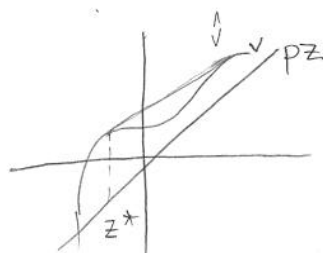
$v^*(p)$ is convex

$\Rightarrow v^*(\lambda p + (1-\lambda)p') \leq (1-\lambda)v^*(p) + \lambda v^*(p')$ } confirm this!

Relationship b/w $v^*(p)$ and $(\hat{v} = \text{conc } v)^*(p)$?

$v^*(p) = \hat{v}^*(p)$ by proposition: If $p \in \partial v(z)$, then $v(z) = \hat{v}(z)$

$$v^*(p) = \sup_z \{v(z) - pz\}$$



v^* maximizes this positive difference

This will still be the z^* of concv

Application to a PTE:

We cannot separate out E from \mathcal{E}

$$I = \{1, \dots, n\}$$

v_i completely characterizes i 's preferences.

$$v_i: \mathbb{R}^l \rightarrow [-\infty, \infty)$$

Cannot disconnect v -fn from endowments since endowments are built into the v -fns.

$$\mathcal{E} \rightarrow v = (v_1, \dots, v_n)$$

$$[(z_i, m_i)] = [(z_1, m_1), \dots, (z_n, m_n)] - \text{allocations}$$

$$(p, 1) - \text{prices}$$

Such that

Utility is maximized (equivalent ways of writing this)

1) $(z_i, m_i) \in d(v_i, p, 0) \quad \forall i$

2) $p \in \partial v_i(z_i) \quad \forall i$

3) $v_i(z_i) - p z_i = v_i^*(p)$

The trade is attainable:

$$\sum_{i=1}^n (z_i, m_i) = (0, 0) \in \mathbb{R}^{2+1}$$

Is this equilibrium efficient?

An allocation that is feasible and maximizes the sum of utilities is efficient under QL preferences

$$\max \{ \sum v_i(z_i) \mid \sum z_i = 0 \}$$

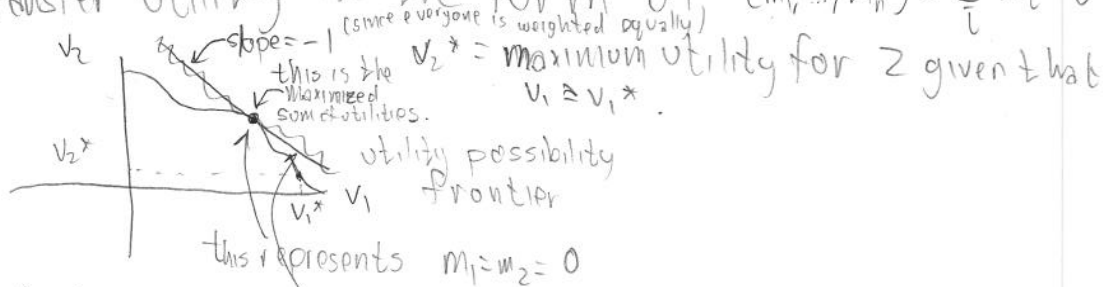
$$v_i(\lambda z + (1-\lambda)z') > \lambda v(z) + (1-\lambda)v(z'), \quad z \neq z'$$

↳ strict concavity \Rightarrow uniqueness

1st step] Maximize the sum of utility

2nd step] Transfer utility in the form of $(m_1, \dots, m_n) \ni \sum m_i = 0$

let $n=2$



$$\max_{z_1} v_1(z_1) + v_2(-z_1)$$

this becomes the Pareto frontier